

## Type I/II sums and sieves

Motivating Q: Given  $A \subseteq [x, 2x]$ , how many primes are in  $A$ ?

Examples: 1)  $A = \{p+2 \in [x, 2x] : p \text{ prime}\} \leftrightarrow$  Twin primes

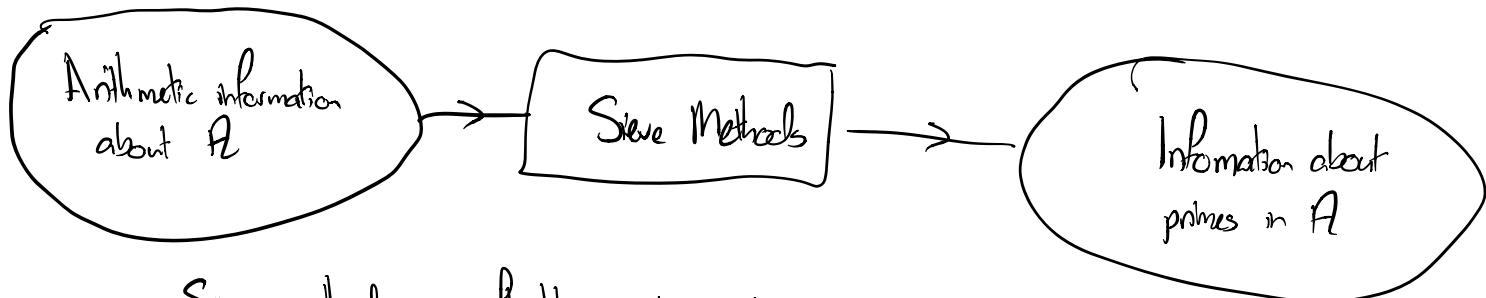
2)  $A = \{n^2 + 1 : n \in [\sqrt{x}, (2x)^{1/2}]\} \leftrightarrow$  Primes of the form  $n^2 + 1$ .

3)  $A = \{n \in [x, 2x] : \|n\alpha\| \leq x^{0.1}\} \leftrightarrow$  Prime in Bohr sets

(here  $\alpha$  is a fixed irrational,  $\|\cdot\| =$  distance to the nearest integer)

Expect asymptotic formulae in all cases, but only known for ③.

In general this question is v. difficult! But Type I/II sums + sieve methods provide a general framework for approaching these problems.



Sieve methods are flexible combinatorial techniques to turn arithmetic information that we can prove into arithmetic information about primes in  $A$ .

Q: Given  $A$  of interest, what arithmetic information can we prove about  $A$ ?

Q: How would we turn this into information about primes in  $A$ ?

Basic/traditional sieve methods just use 'Type I' arithmetic information.



Distribution of  $A$  is residue classes

Notation:  $A_d := \{a \in A : d \mid a\}$ .

Example ①  $A = \{p+2 \in [x, 2x] : p \text{ prime}\}$

$$\text{Bombieri-Vinogradov} \Rightarrow \sum_{d \leq x^{\frac{1}{2}-\varepsilon}} \left| \#A_d - \frac{g(d)}{d} \#A \right| \ll_{A,\varepsilon} \frac{\#A}{(\log x)^A}.$$

↑ where  $g(d) = \begin{cases} \frac{d}{\phi(d)}, & \text{if } d \text{ odd,} \\ 0, & \text{if } d \text{ even.} \end{cases}$

*Conjectured that we can set some result for  $d \leq x^{1-\varepsilon}$ .*

$$\textcircled{2} A = \{ n^2 + 1 \leq [x, 2x] \}$$

$$\sum_{d \leq x^{\frac{1}{2}-\varepsilon}} \left| \#A_d - \frac{g(d)}{d} \#A \right| \ll_{A,\varepsilon} \frac{\#A}{(\log x)^A}.$$

↑ where  $g(d) = \#\{ m \in \mathbb{Z}/d\mathbb{Z} : m^2 \equiv -1 \pmod{d} \}$

This is essentially best possible!

$$\textcircled{3} A = \{ n \in [x, 2x] : \| \alpha n \| \leq x^{0.1} \}$$

Can get reasonable understanding of  $A$  in arithmetic progressions here too.

In these setups this arithmetic info gives satisfying answer by studying the number of elements of  $A$  with no small prime factors.

$$S(A, z) = \#\{ a \in A : a \text{ has no prime factors} \leq z \}.$$

$$\text{e.g. } S(A, (2x)^{\frac{1}{2}}) = \#\{ \text{primes in } A \}.$$

$$\begin{aligned} S(A, x^{\frac{1}{100}}) &= \#\{ \text{elements of } A \text{ with no factors} \leq x^{\frac{1}{100}} \} \\ &\leq \#\{ \text{elements of } A \text{ with at most 101 prime factors} \}, \end{aligned}$$

$$S(A, 6) = \#A - \#A_2 - \#A_3 - \#A_5 + \#A_6 + \#A_{10} + \#A_{15} - \#A_{30}$$

$\therefore S(A, z)$  is completely determined by  $\#A_d$  for  $d \mid \prod_{p \leq z} p$

But / This involves  $d$  which are very big when  $z \geq \log x$ .

Theorem (Linear sieve). Let  $A \subseteq [x, 2x]$  such that

$$\#A = 1 - \#A$$

$$\sum_{d \leq x^{\epsilon}} |\#A_d - \frac{g(d)}{d} \#A| \leq \frac{\#A}{(\log x)^{100}}$$

for some multiplicative function  $g(d)$  with ①  $\sum_{p \leq y} \frac{g(p)}{p} = \sum_{p \leq y} \frac{1}{p} + O(1)$  ← " $g(p)$  is 1 on average"

$$\textcircled{B} \quad g(p) \leq (1-\varepsilon)p \quad \leftarrow \begin{matrix} \text{"not all elements} \\ \text{are a multiple of } p \end{matrix}$$

Then

$$S(A, x^c)$$

$$\mathcal{G}\left(S\left(\frac{1}{c}\right) + o(1)\right) \frac{\#A}{c \log x} \leq \#\{\text{elements of } A \text{ with no factors } \leq x^c \leq S\left(\frac{1}{c}\right) + o(1)\} \frac{\#A}{c \log x} \mathcal{G}$$

for some continuous functions  $\mathcal{G}, F$  with

$$\mathcal{G} = \prod_p \left(1 - \frac{g(p)}{p}\right) \left(1 - \frac{1}{p}\right)^{-1} \quad \begin{aligned} F(s) &= \frac{2}{s} & 1 \leq s \leq 3 \\ \mathcal{G}(s) &= 0 & 0 \leq s \leq 2 \\ \frac{\partial}{\partial s} (sF(s)) &= \mathcal{G}(s-1) & (\text{Fact: } F(s), \mathcal{G}(s) \rightarrow e^{-s}) \\ \frac{\partial}{\partial s} (s \mathcal{G}(s)) &= F(s-1) & \text{as } s \rightarrow \infty \end{aligned}$$

Amazing fact: This is essentially best possible!

$$\exists \text{ sets } A^+ \text{ and } A^- \text{ s.t. } \sum_{d \leq x^{1-\epsilon}} \left| A_d^+ - \frac{g(d)}{d} \#A^+ \right| \ll_A \frac{\#A^+}{(\log x)^A}$$

$$\sum_{d \leq x^{1-\epsilon}} \left| A_d^- - \frac{g(d)}{d} \#A^- \right| \ll_A \frac{\#A^-}{(\log x)^A}$$

$$\left( S\left(\frac{1}{c}\right) + o(1) \right) \frac{\#A^-}{c \log x} = S(A^-, x^c) \leq S(A^+, x^c) = \left( F\left(\frac{1}{c}\right) + o(1) \right) \frac{\#A^+}{c \log x}$$

for all  $c \in (0, 1)$ .